

Verification of Model Transformation using Alloy

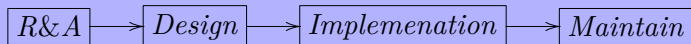
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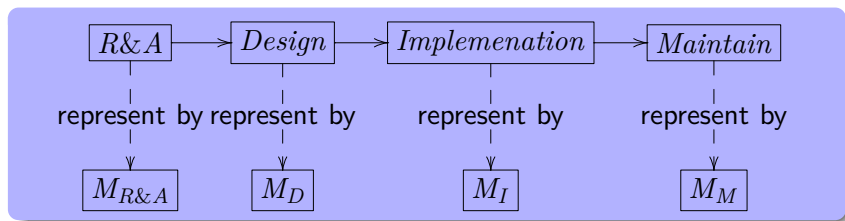
May 30, 2014

- 1 Introduction and Motivation
- 2 Graph-based Model Transformation System
- 3 Encoding of graph-based Model Transformation
- 4 Verification
- 5 Conclusion

Model-Driven Engineering (MDE)



Model-Driven Engineering (MDE)

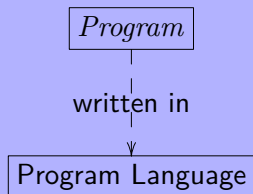


- In model-driven engineering, models are
 - Primary artifacts
 - Used to specify, generate and maintain code

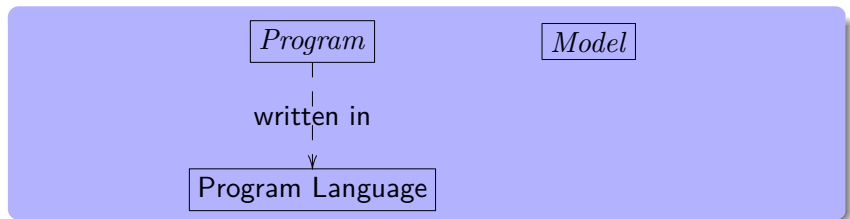
Model Specification and Conformance

Program

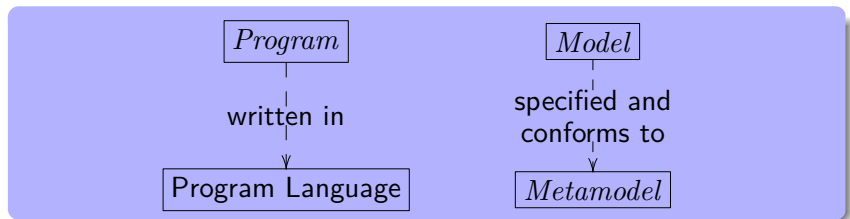
Model Specification and Conformance



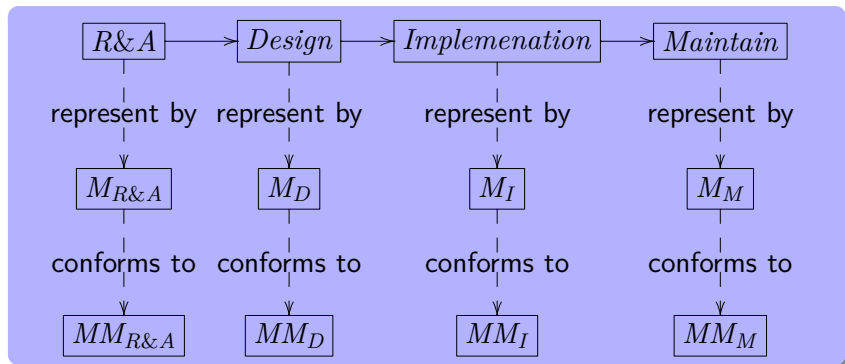
Model Specification and Conformance



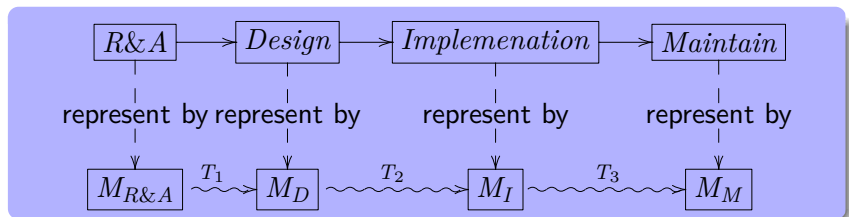
Model Specification and Conformance



Model Specification and Conformance

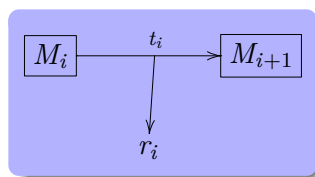
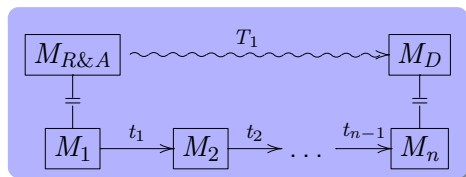


Model Transformation



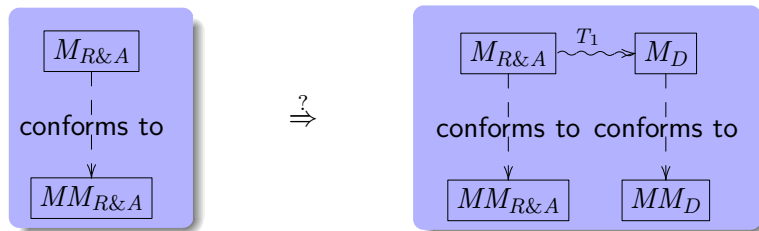
- In model-driven engineering, models are
 - Primary artifacts
 - Used to specify, generate and maintain code
 - Manipulated by model transformations

Model Transformation

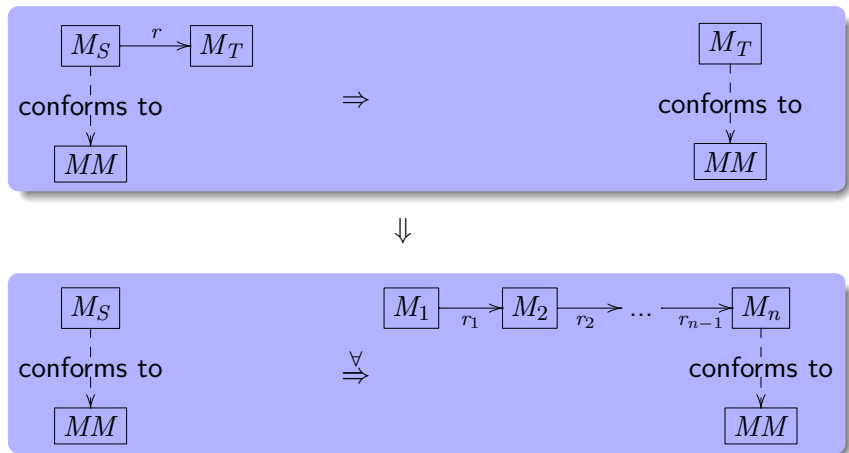


- Model transformations are executed by applying model transformation rules.
- Model transformation rules
 - are defined on metamodel level
 - tell how to execute a transformation, i.e., generate a target model from a source model

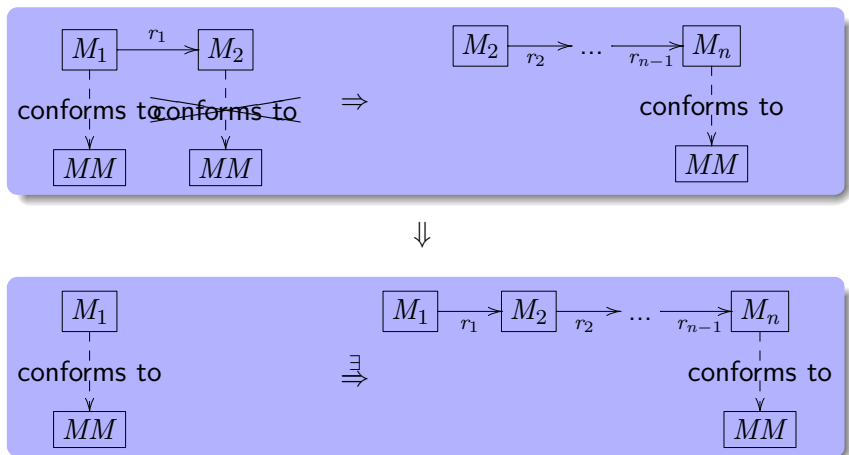
Model Transformation Conformance Problem



Direct Condition



Sequential Condition



Bounded Verification Approach

- Model Transformation System
 - Metamodel
 - Model Transformation Rules
- Conditions
 - Direct Condition
 - Sequential Condition

encoded to

Alloy Specification

checked by

Alloy Analyser

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UML+OCL DPF	Structure diagram diagram	Constraints text diagram
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- A fully diagrammatic specification framework for MDE
- Aims to be a diagrammatic formalism to specify and reason about models and model transformation

- Dijkstra's algorithm ensures that a critical resource is exclusively accessed by one process at a time
- When a process needs to access a resource, the process first sends a request. After the process has acquired the turn, it is allowed to access the resource only if no other process competes with it.

Diagram Predicate Framework

Models are specified by a diagrammatic specification with a graph structure

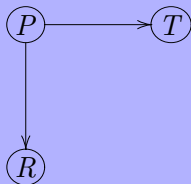


Diagram Predicate Framework

Constraints are added on part of the graph structure

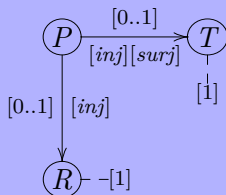
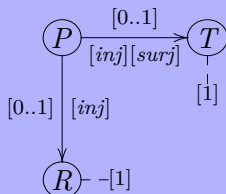
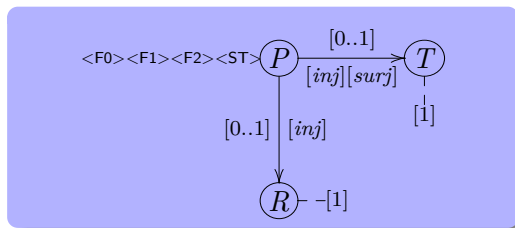


Diagram Predicate Framework



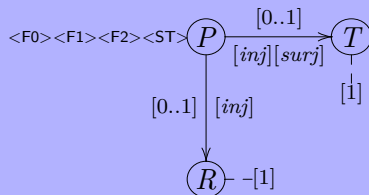
Π^{Σ_1}	α^{Σ_1}	Proposed Vis.	Semantic Interpretation
[m..n]	$1 \xrightarrow{a} 2$	$X \xrightarrow[f]{[m..n]} Y$	$\forall x \in X : m \leq f(x) \leq n, \text{ with } 0 \leq m \leq n \text{ and } 1 \leq n$
[inj]	$1 \xrightarrow{a} 2$	$X \xrightarrow[f]{[inj]} Y$	$\forall x_1, x_2 \in X : f(x_1) = f(x_2) \implies x_1 = x_2$

Diagram Predicate Framework



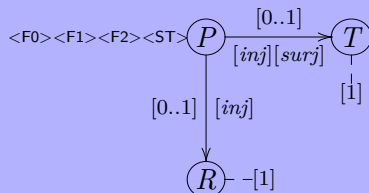
$\Pi^{\Sigma^2_2}$	$\alpha^{\Sigma^2_2}$	Proposed Vis.	Semantic Interpretation
[F0]	1	X	$\langle F0 \rangle$
[F1]	1	X	$\langle F1 \rangle$
[F2]	1	X	$\langle F2 \rangle$
[ST]	1	X	$\langle ST \rangle$

Constraints

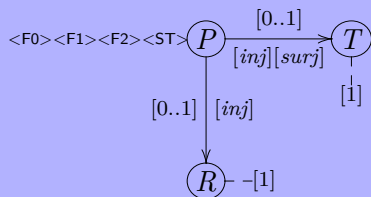


- 1 There is exactly one resource in a model, which is ensured by the multiplicity $[1]$ on R . A similar constraint also applies to T
- 2 There is exactly one process owning the turn, which is ensured by the multiplicity $[0..1]$, $[inj]$ and $[surj]$ on $P \rightarrow T$

Constraints

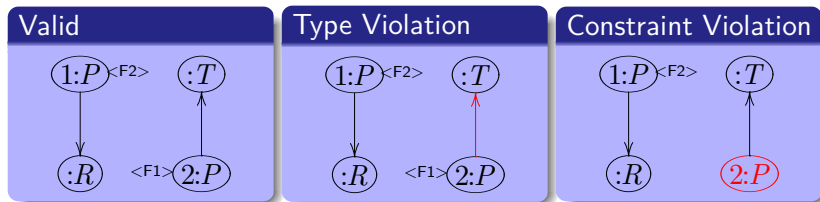


- 3 Each process may access the resource, represented by the arrow $P \rightarrow R$. A multiplicity $[0..1]$ is used to restrict only one such arrow for each process
- 4 In the system, at most one process is accessing the resource, ensured by $[0..1]$ and $[inj]$ on the arrow $P \rightarrow R$



- 5** Each process should have exactly one of the flags (or states) $F0$, $F1$, $F2$ or ST (explained below). These flags are specified as annotations $[F0]$, $[F1]$, $[F2]$, $[ST]$ marking the processes (instances of P)

Conformance Example



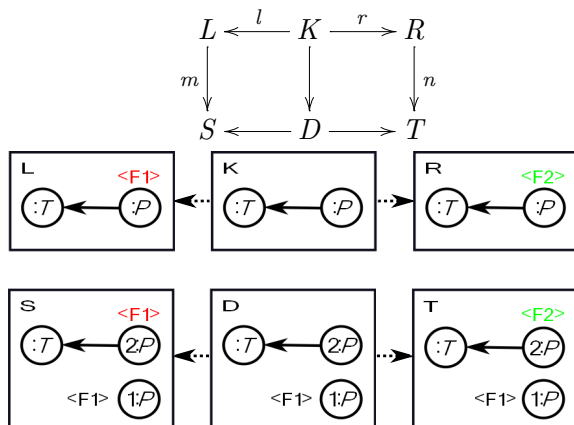
Model Transformation

- Model Transformation is executed according to model transformation rules.
- Model transformation rules specify how to transform a model.
- Graph-based model transformation can be present as graph production.
- Each graph production $L \xleftarrow{l} K \xrightarrow{r} R$ consists the left-hand side, the gluing graph and the right-hand side, where graph morphisms l and r are injective.

Model Transformation Rules

Rule	L	K	R
<i>Request</i>	$(:P) \leftarrow \langle F0 \rangle$	$(:P)$	$(:P) \leftarrow \langle F1 \rangle$
<i>SetFlag2</i>	$(:T) \leftarrow (1:P) \leftarrow \langle F1 \rangle$	$(:T) \leftarrow (1:P)$	$(:T) \leftarrow (1:P) \leftarrow \langle F2 \rangle$
<i>CheckTurn</i>	$(1:P) \leftarrow \langle F1 \rangle$	$(1:P)$	$(1:P) \leftarrow \langle ST \rangle$
	$(:T) \leftarrow (2:P) \leftarrow \langle F0 \rangle$	$(:T) \leftarrow (2:P) \leftarrow \langle F0 \rangle$	$(:T) \leftarrow (2:P) \leftarrow \langle F0 \rangle$
<i>GetTurn</i>	$(1:P) \leftarrow \langle ST \rangle$	$(1:P)$	$(:T) \leftarrow (1:P) \leftarrow \langle F2 \rangle$
	$(:T) \leftarrow (2:P)$	$(2:P)$	$(2:P)$
<i>Compete</i>	$(R) \leftarrow (1:P) \leftarrow \langle F2 \rangle$	$(1:P)$	$(1:P) \leftarrow \langle F1 \rangle$
	$(2:P) \leftarrow \langle F2 \rangle$	$(2:P) \leftarrow \langle F2 \rangle$	$(2:P) \leftarrow \langle F2 \rangle$
<i>Access</i>	$(:R) \leftarrow (1:P) \leftarrow \langle F2 \rangle$	$(:R) \leftarrow (1:P) \leftarrow \langle F2 \rangle$	$(:R) \leftarrow (1:P) \leftarrow \langle F2 \rangle$
	$(2:P) \leftarrow \langle F2 \rangle$		
<i>Exit</i>	$(:R) \leftarrow (:P) \leftarrow \langle F2 \rangle$	$(:R) \leftarrow (:P)$	$(:R) \leftarrow (:P) \leftarrow \langle F0 \rangle$

Double-Pushout Approach



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Bounded Verification Approach

- Model Transformation System
 - Metamodel
 - Model Transformation Rules
- Conditions
 - Direct Condition
 - Sequential Condition

encoded to

Alloy Specification

checked by

Alloy Analyser

Alloy consists of

- the Alloy specification language
 - is a declarative language based on relational logic
 - suited to describe complex model structures and constraints
- the Alloy analyser analyses specifications.

- A model structure is defined as *signatures*
- Each *signature* defines a typed element in the structure, representing a set of instances of this type
- Relations among the typed elements are defined by the fields of the signatures
- Constraints on the structure can be defined as *facts* while predicates as *pred*

- find valid instances well-typed by the structure and satisfying its constraints by executing the *run* command
- verify some properties by calling the *check* command to find counterexamples
- NB! Alloy performs a bounded check, i.e., for each signature, a user-defined scope bounds the number of its instances.

Model Transformations Encoding

	Element	Encoded As
General	Model	<i>sig</i> Graph
	Direct Model Transformation	<i>sig</i> Trans
	Model Transformations	<i>sig</i> Path
Specific	Node	<i>sig</i> Node
	Edge	<i>sig</i> Edge
	rule application	<i>pred</i> rule
	constraints	<i>fact</i>

General Part

Model encoding

Models are encoded as the signature *Graph*

```
1 sig Graph{  
2 nodes:set Nodei+...+Nodem,  
3 edges:set Edgej+...+Edgen}
```

```
1 sig Graph{  
2 nodes:set NP+NR+NT+Int,  
3 edges:set APT+APR+APA}
```

```
1 fact{all g:Graph|all n:g.nodes|all e:Edgej+...+Edgen|(e.src=n  
or e.trg=n) implies e in g.edges}
```

General Part

Transformation Encoding

Direct model transformations are encoded as the signature *Trans*

```
1 sig Trans{
2 rule: one Rule,
3 source: one Graph,
4 target: one Graph,
5 dnodes, anodes: set Nodei+...
  +Nodem,
6 dedges, aedges: set Edgej+...
  +Edgen}
```

```
1 sig Trans{
2 rule: one Rule,
3 source: one Graph,
4 target: one Graph,
5 dnodes, anodes: set NP+NR+NT+
  Int,
6 dedges, aedges: set APT+APR+
  APA}
```

General Part

Valid Transformations

```
1 pred valid_trans [t:Trans]{
2 all edge : Edge|(edge.src in t.dnodes or edge.trg in t.
   dnodes) implies edge in t.dedges
3 all edge : Edge|(edge.src in t.anodes or edge.trg in t.
   anodes) implies edge in t.aedges
4 t.dnodes in t.source.nodes and t.dedges in t.source.edges
5 t.anodes in t.target.nodes and t.aedges in t.target.edges
6 t.source.nodes-t.dnodes=t.target.nodes-t.anodes
7 t.source.edges-t.dedges=t.target.edges-t.aedges
8 rule1[t] or ... or rulen[t]}
9 fact{all t:Trans|valid_trans [t]}
```

General Part

A sequence of Transformations

A sequence of direct model transformations is encoded as path signature of length n .

```
1 sig Path3{trans1, trans2, trans3: one Trans}
2 fact{all p:Path3|p.trans1.target=p.trans2.source and p.
   trans2.target= p.trans3.source}
3 fact{Path3.trans1+Path3.trans2+Path3.trans3=Trans}
```


Specific Part

Nodes and Edges

Assuming a metamodel consisting of m nodes and n edges. Nodes and edges are discriminable from name. Each node(edge) named $i(j)$ is encoded as a $Node_i(Edge_j)$ signature.

```
1 sig Nodei{>//1≤i≤m
2 sig Edgej{>//1≤j≤n
3   src:one Nodes//1≤s≤m
4   trg:one Nodet//1≤t≤m
```

Specific Part

Metamodel constraints encoding

Semantics of DPF predicates are specified in Java or OCL. Those can be encoded to FOL, expressed in Alloy.

```
1 pred source_valid[t:Trans]{
2   //multiplicity on T min:1;max:1
3   let count=#NT&t.source.nodes|count>=1 and count<=1
4   //multiplicity on R min:1;max:1
5   let count=#NR&t.source.nodes|count>=1 and count<=1
6   //injective on PT:P->T
7   all n:(NP&t.source.nodes)|no e1,e2:(APT&t.source.arrows)
      |(e1!=e2 and e1.src=n and e2.src=n and e1.trg=e2.trg
      )
8   //injective on PR:P->R
9   all n:(NP&t.source.nodes)|no e1,e2:(APR&t.source.arrows)
      |(e1!=e2 and e1.src=n and e2.src=n and e1.trg=e2.trg
      )
10  ...
11 }
```

Direct Model Transformation Encoding

Matching

According to DPO, in every direct model transformation $t : S \xrightarrow{p} T$ applying rule $p = L \xleftarrow{l} K \xrightarrow{r} R$,

- the source(target) model has one match of the left(right) part of the rule p . The matched parts are deleted or added according to the rule.

```
1 one m:L→S | (all n:m(l(K_N)) | n in t.source.nodes - t.dnodes) and
2   (all e:m(l(K_E)) | e in t.source.edges - t.dedges) and
3   (all n:m(L_N) \ m(l(K_N)) | n in t.dnodes) and
4   (all e:m(L_E) \ m(l(K_E)) | e in t.dedges)
5 one n:R→T | (all n:n(r(K_N)) | n in t.target.nodes - t.anodes) and
6   (all e:n(r(K_E)) | e in t.target.edges - t.aedges) and
7   (all n:n(R_N) \ n(r(K_N)) | n in t.anodes) and
8   (all e:n(R_E) \ n(r(K_E)) | e in t.aedges)
```

Direct Model Transformation Encoding

Exactly one rule is applied

- No elements other than the matched part are deleted or added in the transformation.
 - 1 If more than one elements (e_1, e_2, \dots, e_m) of the same type t are deleted(added)

```
1 all e:t|e in t.dnodes implies (ori=1i=m (  
2   one m:replace(L,e,ei)→S | (all n:m(l(KN))| n in t.  
   source.nodes-t.dnodes) and  
3   (all e:m(l(KE))| e in t.source.edges-t.dedges) and  
4   (all n:m(replace(LN,e,ei))\m(l(KN))|n in t.dnodes) and  
5   (all e:m(replace(LE,e,ei))\m(l(KE))|e in t.dedges)))
```

Direct Model Transformation Encoding

Exactly one rule is applied

- No elements other than the matched part are deleted or added in the transformation.
 - 2 If one or no element of the type t is deleted(added)

```
1 #{trans.dnodes&t}=1          no trans.dnodes&t
2 #{trans.dedges&t}=1         no trans.dedges&t
3 #{trans.anodes&t}=1         no trans.anodes&t
4 #{trans.aedges&t}=1        no trans.aedges&t
```

Direct Model Transformation Encoding

- No elements other than the matched part are deleted or added in the transformation.
 - 3 For each kept node n , if there are both deleted edges(d_1, d_2, \dots, d_m) and added edges(a_1, a_2, \dots, a_n) related to it, its context should be changed as the rule specified. st is *src* or *trg* field of the node depending on how those edges are related to the node.

```
1 one d1, d2, ..., dm, a1, a2, ..., an | d1.st=d2.st=...=dm.st=a1.st=a2.st,
   ...=an.st
```

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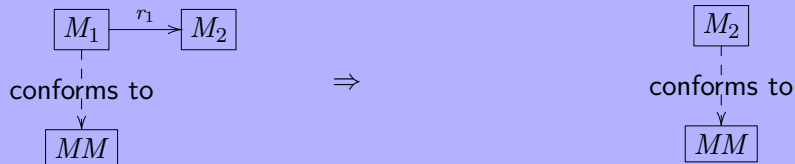
Alloy Specification

checked by

Alloy Analyser

- Verify systems by finding counterexamples
- Verify constraint one by one

Direct Condition



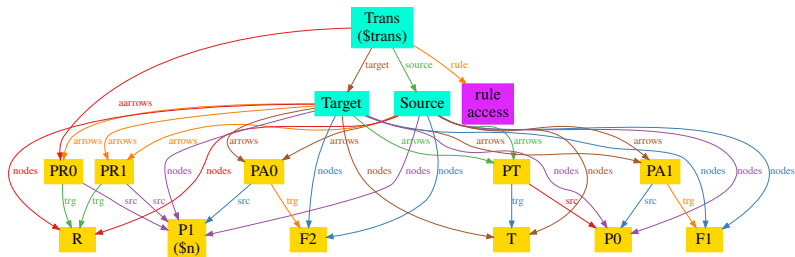
$$\forall S, T : S \triangleright SM \wedge (S \rightarrow T) \Rightarrow T \triangleright TM \quad (1)$$

Commands

- 1 `check{all trans:Trans|all n:(NP&trans.target.nodes)| no e1, e2:(APR&trans.target.arrows)| (e1!=e2 and e1.src=n and e2.src=n and e1.trg=e2.trg)} for 25 but exactly 1 Trans, exactly 2 Graph, exactly 1 Rule`
- 2 `check{all trans:Trans|all n:(NP&trans.target.nodes)| lone e:(APT&trans.target.arrows)|e.src=n} for 25 but exactly 1 Trans, exactly 2 Graph, exactly 1 Rule`
- 3 `check{all trans:Trans|all n:(NF2&tran.target.nodes)|one e: AF2R&trans.target.arrows|e.src=n}for 4 but exactly 1 NR, exactly 1 NT, exactly 1 ATR, exactly 1 Trans, exactly 2 Graph, exactly 1 Rule`

One Counterexample

It shows that 3 constraints are violated. One example is given as following:



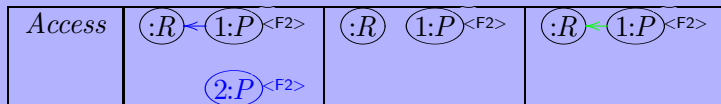
Fix the problem

■ Add more constraints on the metamodel

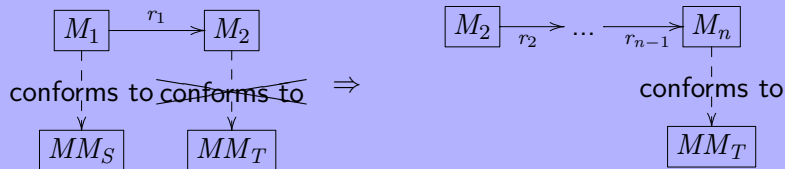
- 1 If a process is accessing the resource, it is at state $\langle F2 \rangle$
- 2 If a process is at state $\langle ST \rangle$, it doesn't own turn

■ Add NACS on certain rules

Revised Rule Access



Sequential Condition



$$\forall \mathbf{S}, \mathbf{T}_0 : (\mathbf{S} \triangleright \mathbf{M} \wedge \mathbf{S} \rightarrow \mathbf{T}_0 \wedge \neg(\mathbf{T}_0 \triangleright \mathbf{M})) \Rightarrow \quad (2)$$
$$\exists \mathbf{T}_1, \dots, \mathbf{T}_n : \bigwedge_{i=1}^{n-1} \neg(\mathbf{T}_i \triangleright \mathbf{M}) \wedge \mathbf{T}_n \triangleright \mathbf{M} \wedge \bigwedge_{i=0}^{n-1} \mathbf{T}_i \rightarrow \mathbf{T}_{i+1}$$

- Verify systems by finding counterexamples
- Verify constraint which violates the direct condition
- Verify stepwise, 2, 3, \dots , n

Object Oriented systems into entity-relationship models

- each class should be transformed into a table
- each table should have a primary key

Path of length 2

```
1 sig Path{trans0,trans1:Trans}{trans0.target=trans1.source}
2 fact{all table:Table&trans0.source.nodes|some pk:PK&trans0.
  source.edges|pk.src=table}
3 fact{some table:Table&trans0.target.nodes|(no pk:PK&trans0.
  target.edges|pk.src=table) and (some pk1:PK&trans1.
  aedges|pk1.src=table)}
```


- Liveness property specify something good will eventually happen

Example

If a process sends request to access the resource(*pre*), it will eventually accesses the resource(*live*).

- Given a liveness property p , if for each possible sequence of transformations $M_0 \rightarrow M_1 \cdots \rightarrow M_k$, some model $M_i (i \in \{0..k\})$ satisfies p , the system satisfies the liveness property p
- a loop sequence is a sequence of transformations $M_0 \rightarrow M_1 \cdots \rightarrow M_k$ where M_k equals to one of the preceding models $M_i (i \in \{0..k\})$ and no model on the sequence satisfies the property

Liveness Properties Result

Loop sequence is found on a path of length 3

```
1 run{some path:Path3|some n:P|pre[path.t1.source, n] and (not
   live[path.t1.target, n]) and (not live[pat4h.t2.target,
   n]) and (not live[path.t3.target, n])}
2 for 20 but exactly 1 Path3 exactly 3 Trans, exactly 3 Graph,
   exactly 1 Rule
3 pred pre[graph:Graph, n:P]{some e:(APA&graph.arrows)|e.src=n
   and e.trg=F1}
4 pred live[graph:Graph, n:P]{(some e:(APA&graph.arrows)|e.src
   =n and e.trg=F2) and (some e:(APR&graph.arrows)|e.src=n)
   }
```

Analysis and Future work

- The approach can be applied to model transformation systems which transforms between different metamodels
- The approach is bounded in that it verifies in a user-define scope.
- The approach is bounded and restricted by scalability problem. We use annotation modelling approach to solve the problem. The result shows the scope is leveled up from 3 to 25. We should find other mechanism to fix the problem, maybe other solver like SMT.

Thank you!

Questions?

For more information visit: <http://dpf.hib.no/>